

Concurrent Binary Search Trees

Supporting Split and Join

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Motivation, Contribution

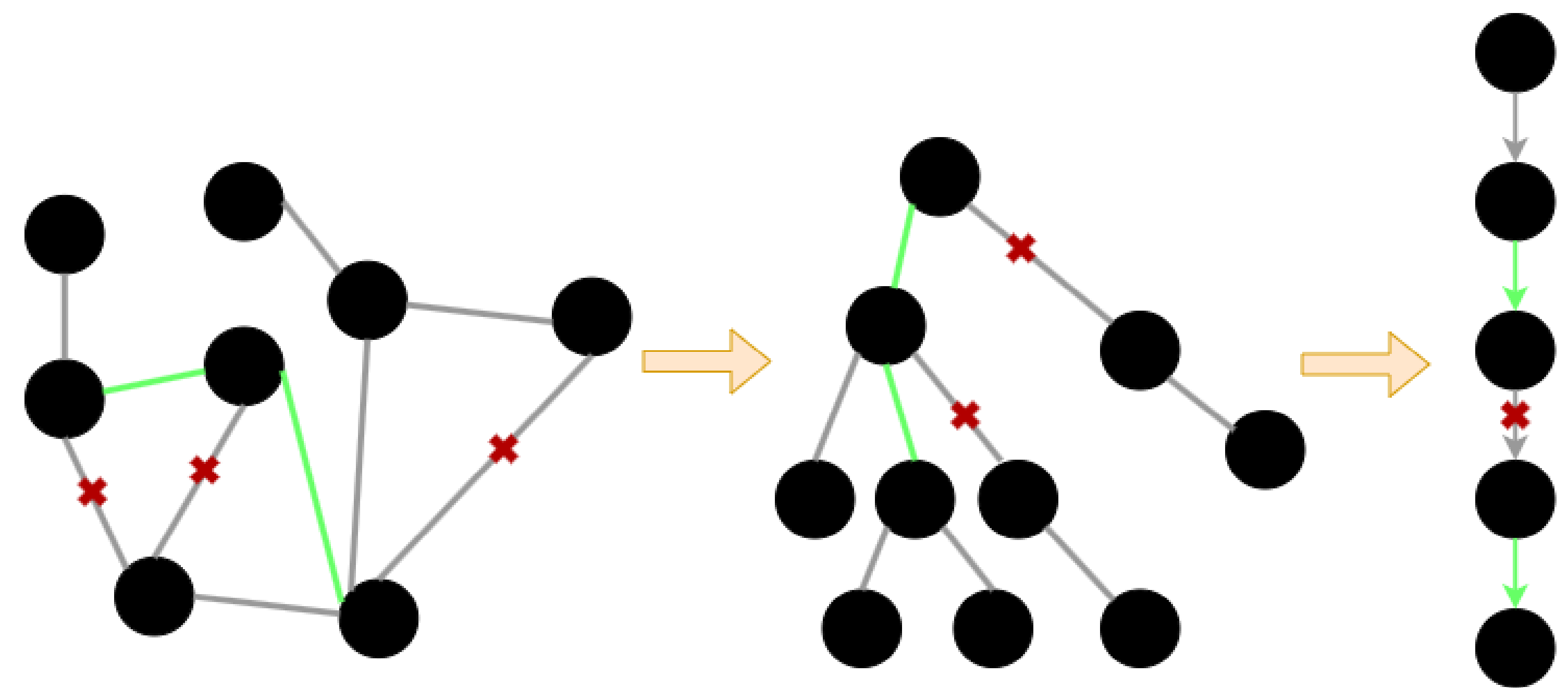
Need for BSTs that support Split and Join Operations

1. Dynamic Graph problems

- Maintain a graph under edge insertions/deletions
- Answer queries efficiently

2. Sequences that support Split and Join

- Essential for some Dynamic Tree data structure (e.g. Link/Cut Trees, Euler Tour Trees)
- Represented using BSTs



BST Operations

□ Split(v): return two tree roots (v : node)

- R_1 : keys $< \text{Key}(v)$
- R_2 : keys $\geq \text{Key}(v)$

□ Join(R_1, R_2): join two tree roots

- Resulting tree has keys of both
- Precondition: the keys of Tree(R_1) are less than the keys of Tree(R_2)

Contribution

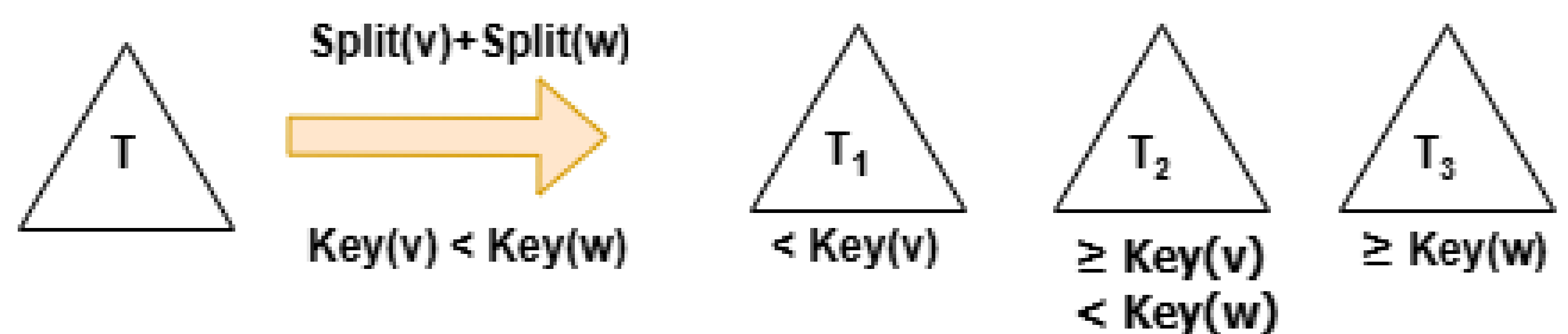
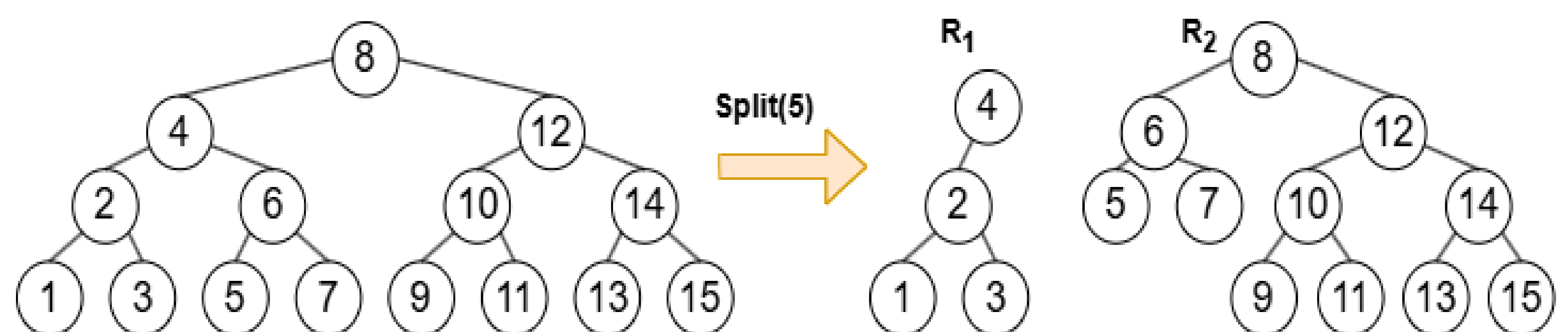
Introducing a lock-based approach for Concurrent Binary Search Trees

- Maintained under Split/Join Operations
- Efficient design: Split operations can combine as they traverse up the Tree
- Intermediate step for solving Dynamic Graph problems in a concurrent setting

Serial Model

Basic Idea: Split(v)

1. v : leaf node (leaf-oriented tree)
2. Keep track of two roots (R_1, R_2) for the constructed trees
3. Initially $R_1 := \text{Null}$, $R_2 := v$
4. Traverse path from leaf v to root
5. If current node is a left child of its parent p :
 - i. Connect R_2 as left child of p
 - ii. $R_2 := p$
6. Else:
 - i. Connect R_1 as right child of p
 - ii. $R_1 := p$



Concurrent Model

Basic Idea: Split(v)

- Similar algorithm to Serial Model
- Locking full path from v to root (synchronization)
- Combining: 2 Splits meet at node m
 - m is the root of T_2
 - One Split continues:
 - replaces its R_1 or R_2 , with one of the roots from the other Split
 - completes the other Split operation along with its own

